More classes of stuck unknotted hexagons

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Abstract

Consider a hexagonal unknot with edges of fixed length, for which we allow universal joint motions but do not allow edge crossings. We consider the maximum number of embedding classes that any such unknot may have. Until now, five was a lower bound for this number. Here we show that there exists a hexagonal unknot with at least nine embedding classes.

1 Introduction

Consider a hexagon in $\mathbb{R}^3$ with fixed edge lengths [1] $\ell_1, \ldots, \ell_6$, and for which we allow universal joint motions but no edge crossings. In other words, we allow the vertices $v_1, \ldots, v_6$ of the hexagon to move freely as long as the edges do not cross or deform. Following the notation of Cantarella and Johnston [CJ98], we denote the space of this hexagon as $Pol_6(\ell_1, \ldots, \ell_6)$. The general question that we study here is: *how many connected components can Pol_6 have, for a suitable choice of $\ell_1, \ldots, \ell_6$?* It is then natural to ask whether each connected component of Pol_6 corresponds to a separate knot type, or if some knot type (say the unknot, or trivial knot) exists in separate components of space. Millet [Mil94] showed that for regular hexagonal unknots there is only one embedding class. The question remained open until

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[1] We use ℓ_i to denote both the edge and its length

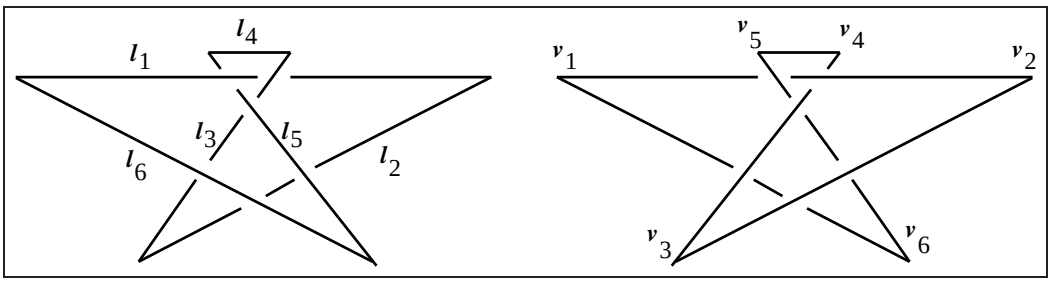

Figure 1: Classes 2 and 3, by Cantarella and Johnston. (not drawn to scale)

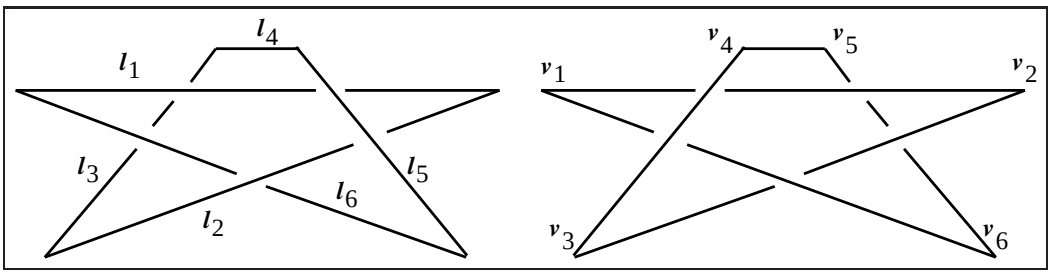

Figure 2: Classes 4 and 5, by Toussaint. (not drawn to scale)

Cantarella and Johnston proved that there exist three connected components in $Pol_6(100, 63.5, 22.7, 5.6, 22.7, 63.5)$, each belonging to the unknot[2]. In other words, they showed that there exist "stuck" unknots which cannot be reconfigured into planar convex embeddings, if the edge lengths are chosen carefully. The first class includes the planar convex embedding. In fact, Aichholzer et al [ADE+01] proved that all planar convex embeddings belong to the same connected component. The remaining two classes, which are left and right hand versions of the same polygon, are shown in Figure 1.

Cantarella and Johnston suspected that these three classes were the only ones for the unknot in Pol_6, but Toussaint [Tou01] showed that there exist two more classes, for $Pol_6(20, 13, 4, 1, 4, 13)$. The two new classes are left and right hand versions of each other, and are shown in Figure 2. In the following section, we show that four more classes can exist for appropriate edge lengths.

[2]The numbers given are approximate

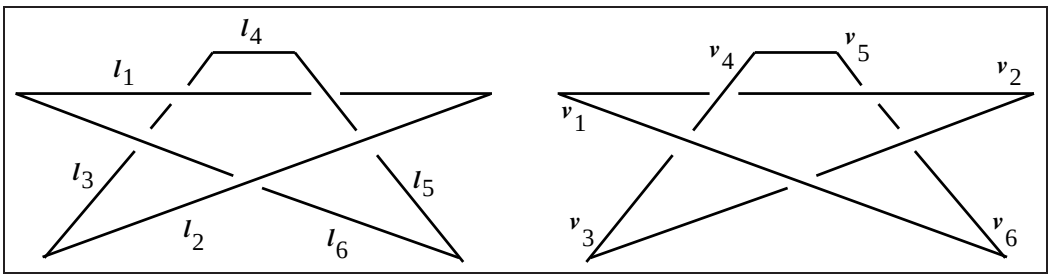

Figure 3: Classes 6 and 7. (not drawn to scale)

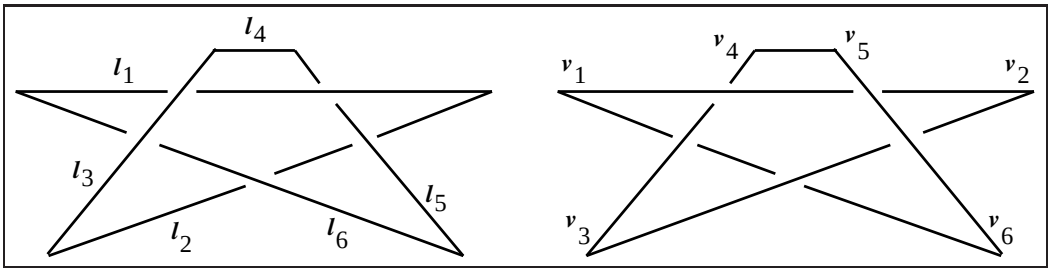

Figure 4: Classes 8 and 9. (not drawn to scale)

2 A hexagonal unknot with nine embedding classes

We consider the space $Pol_6(1, 1, 0.55, \epsilon, 0.55 - \epsilon, 1)$, where $\epsilon \leq 0.01$ (i.e. sufficiently small). Here, we show that this space has at least nine connected components corresponding to the unknot. In other words there are at least nine embedding classes of the unknot for our given edge lengths.

The first class contains all planar convex embeddings. The next four are similar in shape to those of Cantarella and Johnston and those of Toussaint. However, since we are using different edge lengths it is necessary to verify that our classes of hexagons are still in different connected components of space. The four new classes are shown in Figures 3 and 4.

Cantarella and Johnston showed that a sufficient condition for their classes to be in different connected components than the convex class is that $(\ell_3 + \ell_4)^2 < \ell_1\ell_2 - (\ell_1)^2/2$. Our chosen edge lengths satisfy this condition, so we can say that the proof of Cantarella and Johnston holds for our hexagon.

Of course, it is not enough just to show that every class is in a different connected component than that containing the convex class.

Unfortunately, we can not do the same as above and borrow the proof for Toussaint's classes. A sufficient (but not necessary) condition used in this case is that $\ell_3 + \ell_4 + \ell_5 < min\{\ell_2, \ell_6\}$. Our hexagon is modified enough that this condition no longer holds. It seems that we cannot avoid such a modification, in order to create the new classes.

Even though Toussaint's proof cannot be used here, the intuition is quite the same. For each class, there are a couple of motions that must be made in order to convexify the hexagon. We will show that such motions are impossible to make, and conclude that it is impossible not only to convexify a hexagon belonging to each class but also to reconfigure between the non-convex classes.

Let us now fix a coordinate system in order to view all possible motions. Let ℓ_1 be fixed in the xy plane, and ℓ_6 be constrained to the plane. Specifically, v_1 is at $(0,0,0)$, v_2 is at $(1,0,0)$, and v_6 is at $(v_{6x}, v_{6y}, 0)$ where $v_{6x}^2 + v_{6y}^2 = 1$ Our view will be along the normal of the plane. Thus vertex v_6 may move only in a circle about v_1. We choose to focus on class 9, as shown on the right of Figure 4. Intuitively we can see that to convexify this hexagon we would have to pass v_3 over ℓ_6 and into the plane, or we would have to pass ℓ_5 over v_2 so that v_5 could be placed in the plane.

We now proceed to describe certain constraints in the configuration of our polygon. The polygon shown in Figure 4 possesses certain properties, listed below:

1. $v_{5z} > 0$, and $v_{4z} < 0$. Thus we can define the point $P = P_{(x,y)} = (P_x, P_y)$ to be the intersection of ℓ_4 with the plane. This also means that v_4 and v_5 are at a distance of at most ϵ from the plane.

2. $P_y > 0$, and $\epsilon < P_x < 1 - \epsilon$.

3. $v_{3z} > 0$.

We will show that, while reconfiguring the polygon, it is impossible to change any single property without causing an intersection or contradicting the validity of at least one other. We will also show that we cannot change more than one of these properties simultaneously. This implies that the properties are always true. Note that *changing* a property is done via a continuous motion of the polygon. Thus for example, property 1 changes the

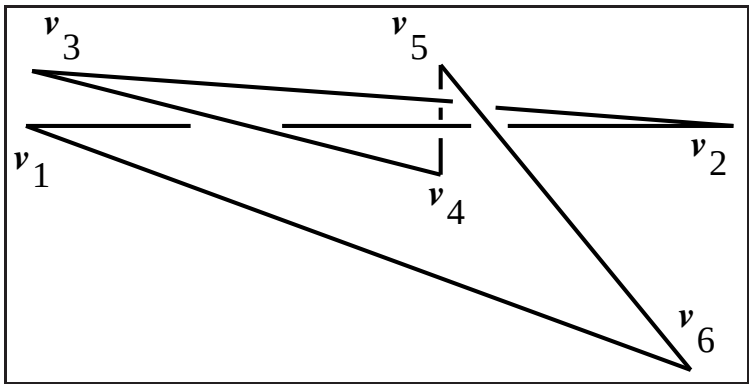

Figure 5: Alternate configuration of class 9. (not drawn to scale)

moment that $v_{5z} = 0$ and/or $v_{4z} = 0$. In case we attempt to set both of these vertices into the plane simultaneously, the point P will be defined to be the last unique point of intersection of ℓ_4 with the plane.

First let's see if we can change only the third property. Consider the case where $v_{3y} < 0$. v_3 can be placed in the plane only if the angle at v_2 or the angle at v_1 opens to more than $\pi/3$. This would mean that the distance from v_3 (or v_6 respectively) to ℓ_1 becomes greater than $\frac{\sqrt{3}}{2}$. The distance to P is at least as great, since $P_y > 0$. However, the distance from v_3 (v_6) to P can be at most the sum of the lengths $\ell_3 + \ell_4$ (or $\ell_5 + \ell_4$ respectively). These sums are less than $\frac{\sqrt{3}}{2}$. Thus, by contradiction, we conclude that v_3 cannot be placed in the plane if $v_{3y} < 0$. Allowing properties 1 and/or 2 to change simultaneously with property 3 does not affect the arguments given above. At the "critical" moment that we attempt to place v_3 into the plane, we still have $P_y \geq 0$, for example.

It remains to be seen if we can change only the third property, when $v_{3y} > 0$. Even reaching a configuration where $v_{3y} > 0$ requires a subtle motion. It can be done by bringing P sufficiently close to ℓ_1, which allows us to obtain $v_{4y} < 0$. Then ℓ_2 can move in a clockwise motion as viewed from the normal to the plane. The final configuration is shown in Figure 5. During the motion described above, we have $0 \leq v_{3x} \leq 1$. Thus the x-coordinates of both endpoints of ℓ_3 are within this range. This means that if ℓ_3 is to avoid intersection with ℓ_1 it is in fact necessary to have $v_{4y} < 0$. Keeping in mind that $v_{4z} < 0$ (property 1), we see that it is impossible to place v_3 in the plane without causing an intersection between ℓ_3 and ℓ_1. We conclude

that v_3 always remains above the plane, as long as properties 1 and 2 are true. Again, allowing properties 1 and/or 2 to change at the same time (i.e. allowing $P_x = \epsilon$ *and/or* $v_{4z} = 0$) doesn't affect our arguments. Thus property 3 cannot change, either on its own, or at the same time as properties 1-2.

We now focus on the second property: To obtain $P_y < 0$, we must *first* have $P_x < 0$ or $P_x > 1$. This is because of the position of ℓ_1. So we would be requiring that $v_{4x} < \epsilon$ or $v_{4x} > 1 - \epsilon$. The same holds for v_{5x}. In other words, v_4 and v_5 would have to approach to within a distance ϵ from the halfline $x = 0, y > 0, z = 0$ or the halfline $x = 1, y > 0, z = 0$. By examining the triangle $v_1v_5v_6$ (which can be made to approach arbitrarily close to the plane if ϵ is small enough) and using elementary trigonometry, we can see that the former case is impossible. On the other hand, for v_4 to approach the halfline $x = 1, y > 0, z = 0$, we would need the angle $v_1v_2v_3$ to reach $\pi/3$. As shown previously, this cannot be done when $v_{3y} < 0$. We also established that if $v_{3y} > 0$, as shown in Figure 5, then $v_{4y} < 0$. This means that the angle $v_1v_2v_3$ cannot open to $\pi/3$, by the same logic used when dealing with changing property 3. Thus we must always have $0 < P_x < 1$ and $P_y > 0$. Once again, the argument given is not affected by allowing property 1 to change simultaneously.

Finally, we can examine the first property: Given property 2, we cannot place v_5 into the plane because this would cause an intersection between the polygonal arc Pv_5v_6 and ℓ_1. Also, we cannot place v_4 into the plane because the arc v_3v_4P would intersect ℓ_1. More specifically, if $v_{3y} \geq 0$, bringing v_4 to the plane would cause an intersection between ℓ_4 and ℓ_1, since $P_y > 0$. If $v_{3y} < 0$, the intersection would be between ℓ_1 and either ℓ_3 or ℓ_4 depending on the value of v_{4y}. We conclude that v_4 (v_5) must always be below (above) the plane.

Having established property 1, we know that class 9 cannot be reconfigured to classes 2,5,7,8. Now consider class 6: ℓ_2 passes above the polygonal arc Pv_5v_6, whereas in class 9 it passes beneath the arc. In both cases, ℓ_2 is above the plane. We know that starting from class 9 we cannot position v_3 into the plane. This implies that if we are to reconfigure between these classes, v_3 must pass under ℓ_5. This cannot be done, due to the angle restrictions mentioned while handling property 1. In fact, if this were possible, then v_3 could be placed in the plane. The same applies for class 3. Finally, the only difference between classes 4 and 9 is in the crossing of ℓ_3 and ℓ_6, as we view along the normal to the plane. To change this crossing, the point Q, defined earlier, would have to move to an intermediate position such that

$Q_y < v_{6y}$. This involves passing v_3 under ℓ_5, which cannot be done, as we have just seen.

We can now say that class 9 cannot be reconfigured to any of the other classes shown. If we ignore the small difference in the lengths of ℓ_3 and ℓ_5, which does not play a role in the proofs given above, then the four new classes are either left/right hand versions or rigid transformations of each other. Thus the arguments given for class 9 hold for classes 6–8. As for our modified examples of Toussaint's classes, the proofs are similar. This leads us to the following theorem:

Theorem 2.1 *For suitable choices of edge lengths, there are at least nine connected components belonging to the unknot in Pol_6.*

3 Remarks

Although we show that there are at least four more embedding classes for Pol_6 we note that our examples are not very "stable", in the sense that the slightest change in the length of either ℓ_2 or ℓ_6 would result in the loss of four classes: those of Toussaint, and two new ones. If we made both lengths larger, then we would only have the three classes of Cantarella and Johnston. An open problem is therefore to find suitable edge lengths so that Pol_6 has more than five "stable" embedding classes.

As was mentioned by Cantarella and Johnston, these results can be extended to Pol_n, by replacing the shortest edge with a chain. Alternatively, a link or a knot can be placed to obtain similar results.

Calvo showed that even if edge lengths can vary, there exist four components of trefoil knots in Pol_6 (two left-handed and two right-handed). In fact, no other type of knot can be formed with only six edges [Cal01]. It was tempting to see whether we could modify our classes, by changing some crossings, to obtain more than four components corresponding to the trefoil knot. The examples of Cantarella and Johnston cannot be transformed into trefoils, and the remaining six non-planar unknots produce only four trefoil classes, due to symmetry conditions. It remains an interesting open problem to determine if more than four trefoil embeddings exist in Pol_6 when edge lengths are fixed.

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