

## Lecture 8

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In this class, we did an experiment using Sudoku to implement Merkle puzzles as a public key distribution scheme. Please see Merkle's account of how his idea was initially rejected [1] and the wiki on Merkle's puzzles [2]. First we consider how to generate a key from a Sudoku game.

## 1 Keys from Sudoku

We illustrate with 4 X 4 Sudoku games for simplicity. Start with such a game with 5 of the 15 squares filled out. We label the other 10 squares with the variables  $k_0, \dots, k_9$  in row by row fashion.

For example:

$k_0$	$k_1$	1	$k_2$
$k_3$	$k_4$	4	3
1	$k_5$	2	$k_6$
$k_7$	2	$k_8$	$k_9$

Then the solution to the game

3	4	1	2
2	1	4	3
1	3	2	4
4	2	3	1

gives us a key  $K = (3, 4, 2, 2, 1, 3, 4, 4, 3, 1)$ . We could use this key in a Vigenère cipher, where each message text letter is shifted in alphabetic order by the appropriate key. So for example the message text  $M=HELLO WORLD$  would be encoded  $C=KINNP XSVOE$ .

We can imagine Alice and Bob both carry a book of Sudoku puzzles and agree in advance on a certain puzzle number  $p$ . Then if Alice wishes to send a message to Bob she can simply solve puzzle  $p$  and extract its key  $K$  as above. She encodes her message with  $K$  and sends it publicly to Bob. Bob can solve also solve puzzle  $p$ , get  $K$  and decode her message. It requires that both Alice and Bob solve a single Sudoku game. An eavesdropper Eve receiving  $C$  would have to brute force solve the puzzles, even if she finds the Sudoku book. If the message  $M$  was completely random Eve

could never be sure when she had decoded C. Of course Alice does not solve the puzzle by writing directly in the book itself - this would give it away!

## 2 Merkle's public key distribution based on Sudoku

So far we have seen that a classical key could be generated from a Sudoku game. We will use Merkle's idea to convert this into a public key distribution scheme between Alice and Bob. First we need to compute an identifier for any given key. One way would be to compute a simple hash function, such as  $3 \sum_{i=0}^{i=4} k_i + \sum_{i=5}^{i=9} k_i$ . In this case the identifier would become  $3*12+15=51$ . We assume that this hash function is publicly known.

The protocol for Alice and Bob to agree on a common password M is as follows. We assume that Alice has generated N Sudoku games and published them in a book. Since she knows all the solutions, she can compute a list containing the identifier of each puzzle. We assume she has also deleted any game that would have created a duplicate identifier.

- Bob privately chooses a random Sudoku game from his copy of the book and solves it. He gets a key K and computes its identifier I.
- Bob publicly announces I.
- Alice looks up the identifier I in her list and finds the key K.
- Alice codes up a password M for Bob using K getting code C. She sends C to Bob publicly.
- Bob uses the key K to recover the password M from C
- Alice and Bob communicate using shared password M

As an example, suppose Bob chose the above Sudoku game, getting  $K = (3, 4, 2, 2, 1, 3, 4, 4, 3, 1)$  and identifier  $I=51$ . He announces  $I=51$  publicly. Alice looks up I in her list and recovers K. She chooses the password  $M=HELLO WORLD$  which when encoded by K gives  $C=KINNP XSVOE$ . Alice announces C publicly. Bob now uses K to recover M from C by shifting the code letters backwards in alphabetic order. Using M as a key they can communicate privately.

## 3 Exercises

Many cryptographic schemes have a "trap door", which is a short-cut method to avoid doing a brute force search. The method above admits several such trap-doors. Consider the list of  $N=8$  puzzles below.

- (a) Explain how to compute the signature of 7 of the 8 puzzles without solving any of the puzzles at all!
- (b) For the remaining puzzle show how to compute its signature by solving for only two empty squares.
- (c) Suggest a stronger method of obtaining a signature, and argue why you need to solve each complete puzzle to compute this signature. Compute the signature for each of the eight puzzles.
- (d) Explain why  $M=HELLO WORLD$  is not a good choice of secret password by Alice. Ie., explain how using simple English text could make Eve's job easier. What is a good choice for Alice?

①

		2	
	1		4
4		1	
	2		

②

	4		
1		2	
	2		3
		4	

③

	2	1	
1			4
	4	3	

④

		4	1
4			
2	3		
			3

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⑤

			4
4	3		
		1	3
3			

⑥

		3	
	3		2
3		2	
	4		

⑦

2			4
		1	
	4		
3			1

⑧

4			3
		2	
	4		
3			2

## References

- [1] Merkle's project proposal: <http://www.merkle.com/1974/>
- [2] Merkle's puzzles: <http://en.wikipedia.org/wiki/Merkle>